

## Notes on Kahan's *A Distillation Program*

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### ABSTRACT

In the 8th set of Kahan course notes, there was a 1-page entry "A Distillation Program" (entry 3.5). My copy was almost illegible, which did not help any in trying to understand it.

The following is a recapitulation of the ideas in that Kahan note. It is more legible, and may or may not be clearer.

*Exact 2-operand addition:* Let  $Z$  be  $p+q$  (machine precision, rounded),  $\text{abs}(p) > \text{abs}(q)$ . Let  $z$  be  $(p-Z)+q$  (operations done at machine precision). Then  $Z+z$  is exactly  $p+q$ ; and  $z$  is as small as possible, among all pairs  $Z+z$  for which this holds.

*Exact N-operand addition:* We try to span the nonzero bits of the exact  $N$ -operand sum with (approximately) as few terms as possible. This is called "distillation". We use two data-structures for this: one ("the list") holds all current terms, in order of absolute value. The other ("the heap") other holds pairs of successive list-entries, sorted in order of  $\text{floor}(\log_2(p)) - \text{floor}(\log_2(q))$ . In fact, all we need in the 2nd data-structure is the heap-property, i.e. the ability to find its smallest member fast.

A balanced tree (2-3 or AVL) will do fine for both data-structures ( $\log N$  cost for any operation).

We can assume all entries are initially nonzero; otherwise, we just delete the zeros.

Repeat

find smallest number in heap. Locate the two terms  $p, q$ , in list that correspond to it.

Convert  $p+q$  into  $Z+z$ , as above.

Delete  $p$  and  $q$  from list (and update the heap accordingly)

Insert  $Z$  into list (updating the heap accordingly)

if  $z \neq 0$ , insert in list (updating the heap accordingly)

Until "convert" step is a no-op.

*Termination:* At every step, the values in the list become "more unequal": their sum remains constant, but the new list is "more exact" in the following sense: there are numbers such that

$$\sum X_i \text{ over } |X_i| > Z$$

becomes more accurate after the transformation than it was before. But the reverse does not occur.

Since this "more accurate" relation is a strict ordering, the iteration must terminate.

Final state: In the final state of the list, it contains  $N$  floating-point numbers, whose bit-positions never overlap (for otherwise,  $p+q \rightarrow Z+z$  would not have been a no-op). This number is thus within one of the optimum.